

# Categories: How I Learned to Stop Worrying and Love Two Sorts

Alessandra Palmigiano

joint work with

W. Conradie, S. Frittella, A. Tzimoulis and N. Wijnberg.

18 August 2016

Wollic 2016, Puebla, Mexico

# Introduction

**Problem:** Understanding relational semantics for lattice-based (e.g. substructural) logics:

- Not 1 set of possible worlds, but 2 domains: states and **co-states** (???)
- Semantic interpretation of  $\Box$  and  $\Diamond$  very different from the usual clauses in Boolean and distributive settings
$$\mathbb{M}, x \succ \Diamond\phi \quad \text{iff} \quad \text{for all } a \in A, \text{ if } \mathbb{M}, a \Vdash \phi, \text{ then } xRa$$
$$\mathbb{M}, a \Vdash \Diamond\phi \quad \text{iff} \quad \text{for all } x \in X, \text{ if } \mathbb{M}, x \succ \Diamond\phi, \text{ then } a \perp x$$
- Satisfaction  $\Vdash$  and co-satisfaction  $\succ$  (???)

**Proposed solution:** Formulas as categories (as understood in business science).

**Advantage:** Surprisingly natural epistemic interpretation!

# Introduction

**Problem:** Understanding relational semantics for lattice-based (e.g. substructural) logics:

- Not 1 set of possible worlds, but 2 domains: states and **co-states** (???)
- Semantic interpretation of  $\Box$  and  $\Diamond$  very different from the usual clauses in Boolean and distributive settings
  - $\mathbb{M}, x \succ \Diamond\phi$  iff **for all**  $a \in A$ , if  $\mathbb{M}, a \Vdash \phi$ , then  $xRa$
  - $\mathbb{M}, a \Vdash \Diamond\phi$  iff **for all**  $x \in X$ , if  $\mathbb{M}, x \succ \Diamond\phi$ , then  $a \perp x$
- Satisfaction  $\Vdash$  and co-satisfaction  $\succ$  (???)

**Proposed solution:** Formulas as categories (as understood in business science).

**Advantage:** Surprisingly natural epistemic interpretation!

## Two-sorted semantics for lattice-based modal logics

**Polarity.**  $\mathbb{P} = (A, X, \perp)$  with  $A$  and  $X$  sets and  $\perp \subseteq A \times X$ .

**Galois connection.**  $(\cdot)^\uparrow : \mathcal{P}A \rightarrow \mathcal{P}X$  and  $(\cdot)^\downarrow : \mathcal{P}X \rightarrow \mathcal{P}A$  s.t. for all  $B \subseteq A$  and  $Y \subseteq X$ ,

- $B^\uparrow := \{x \in X \mid \forall a(a \in B \rightarrow a \perp x)\}$ ,
- $Y^\downarrow := \{a \in A \mid \forall x(x \in Y \rightarrow a \perp x)\}$ .

**Closed sets.**  $B = B^{\uparrow\downarrow}$  and  $Y = Y^{\downarrow\uparrow}$ .

**Lattice of closed sets.** Let  $C(A)$  (resp.  $C(X)$ ) be the closed subsets of  $A$  (resp.  $X$ ).

$$\mathbb{P}^+ = (C(A), \bigcap, \bigvee, \emptyset^{\uparrow\downarrow}, A) \cong^{\partial} (C(X), \bigcap, \bigvee, \emptyset^{\downarrow\uparrow}, X).$$

**Concept lattice of  $\mathbb{P}$ .** Lattice of tuples  $(B, Y)$  s.t.

$$Y = B^\uparrow \quad \text{and} \quad B = Y^\downarrow.$$

# RS-polarities

**RS-Polarity.**  $(A, X, \perp)$  such that

(S1)  $b \neq c$  implies  $\perp[b] \neq \perp[c]$ , for all  $b, c \in A$

(S2)  $y \neq z$  implies  $\perp^{-1}[y] \neq \perp^{-1}[z]$ , for all  $y, z \in X$

(R1)  $a^{\uparrow\downarrow}$  completely join-irreducible in  $\mathbb{P}^+$  for every  $a \in A$

(R2)  $x^{\downarrow\uparrow}$  completely meet-irreducible in  $\mathbb{P}^+$  for every  $x \in X$

**Lattice of closed sets.**

$$\mathbb{P}^+ = (C(A), \cap, \vee, \emptyset, A) \cong^{\partial} (C(X), \cap, \vee, \emptyset, X).$$

**Perfect lattices and RS-polarities.** For any perfect lattice  $\mathbb{L}$ ,  $\mathbb{L}_+ := (J^\infty(\mathbb{L}), M^\infty(\mathbb{L}), \leq)$  is an RS-polarity.

**Theorem.** For any RS-polarity  $\mathbb{P}$  and perfect lattice  $\mathbb{L}$ ,

$$(\mathbb{P}^+)_+ \cong \mathbb{P} \quad \text{and} \quad \mathbb{L} \cong (\mathbb{L}_+)^+$$

# RS-frames

**Language.**  $\mathcal{L} = 0 \mid 1 \mid \vee \mid \wedge \mid \Box$

**RS-frame.**  $\mathbb{F} = (\mathbb{P}, R)$  such that

- $\mathbb{P} = (A, X, \perp)$  is an RS-polarity
- $R \subseteq A \times X$
- $R[b]$  and  $R^{-1}[y]$  are closed sets, for all  $b \in A$  and  $y \in X$ .

**RS-models.**  $\mathbb{M} = (\mathbb{F}, V)$  s.t.

- $\mathbb{F}$  RS-frame
- for all  $p \in \mathbf{AtProp}$ ,

$$V(p) = (V_1(p), V_2(p))$$

with  $V_1(p) = V_2(p)^\downarrow$  and  $V_2(p) = V_1(p)^\uparrow$ .

# Interpretation of lattice-based modal logic on RS-frames

$\mathbb{M}, a \Vdash 0$       never                       $\mathbb{M}, x \succ 0$       always

$\mathbb{M}, a \Vdash 1$       always                       $\mathbb{M}, x \succ 1$       never

$\mathbb{M}, a \Vdash p$     iff     $a \in V_1(p)$                $\mathbb{M}, x \succ p$     iff     $x \in V_2(p)$

$\mathbb{M}, a \Vdash i$     iff     $a \in V_1(i)$                $\mathbb{M}, x \succ i$     iff     $x \in V_2(i)$

$\mathbb{M}, a \Vdash m$     iff     $a \in V_1(m)$                $\mathbb{M}, x \succ m$     iff     $x \in V_2(m)$

$\mathbb{M}, a \Vdash \phi \wedge \psi$     iff     $\mathbb{M}, a \Vdash \phi$  and  $\mathbb{M}, a \Vdash \psi$

$\mathbb{M}, x \succ \phi \wedge \psi$     iff    for all  $a \in A$ , if  $\mathbb{M}, a \Vdash \phi \wedge \psi$ , then  $a \perp x$

$\mathbb{M}, a \Vdash \phi \vee \psi$     iff    for all  $x \in X$ , if  $\mathbb{M}, x \succ \phi \vee \psi$ , then  $a \perp x$

$\mathbb{M}, x \succ \phi \vee \psi$     iff     $\mathbb{M}, x \succ \phi$  and  $\mathbb{M}, x \succ \psi$

$\mathbb{M}, a \Vdash \Box\phi$     iff    for all  $x \in X$ , if  $\mathbb{M}, x \succ \phi$ , then  $aRx$

$\mathbb{M}, x \succ \Box\phi$     iff    for all  $a \in A$ , if  $\mathbb{M}, a \Vdash \Box\phi$ , then  $a \perp x$

# Standard translation

$$ST_a(0) := a \neq a \quad ST_a(1) := a = a \quad ST_a(p) := P_1(a)$$

$$ST_x(0) := x = x \quad ST_x(1) := x \neq x \quad ST_x(p) := P_2(x)$$

$$ST_a(\mathbf{j}) := \forall x[a \perp x \rightarrow j \perp x] \quad ST_x(\mathbf{j}) := j \perp x$$

$$ST_a(\mathbf{m}) := a \perp m \quad ST_x(\mathbf{m}) := \forall a[a \perp m \rightarrow a \perp x]$$

$$ST_a(\phi \vee \psi) := \forall x[ST_x(\phi \vee \psi) \rightarrow a \perp x]$$

$$ST_x(\phi \vee \psi) := ST_x(\phi) \wedge ST_x(\psi)$$

$$ST_a(\phi \wedge \psi) := ST_a(\phi) \wedge ST_a(\psi)$$

$$ST_x(\phi \wedge \psi) := \forall a[ST_a(\phi \wedge \psi) \rightarrow a \perp x]$$

$$ST_a(\Box\phi) := \forall x[ST_x(\phi) \rightarrow aRx]$$

$$ST_x(\Box\phi) := \forall a[ST_a(\Box\phi) \rightarrow a \perp x]$$

## Main property.

$$\mathbb{M} \Vdash \phi \leq \psi \quad \text{iff} \quad \mathbb{M} \models \forall a[ST_a(\phi) \rightarrow ST_a(\psi)]$$



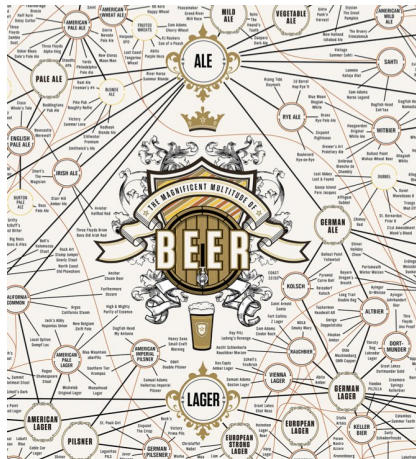
# Categorization theory

From Wikipedia:

*Categorization is the process in which ideas and objects are recognized, differentiated, and understood.*

*Ideally, a category illuminates a relationship between the subjects and objects of knowledge.*

*Categorization is fundamental in language, prediction, inference, decision making and in all kinds of environmental interaction.*



# Categorization theory and RS-models via Formal Concept Analysis

Let  $\mathbb{F} = (\mathbb{P}, R)$  with

- $\mathbb{P} = (A, X, \perp)$  database
- $A$  set of objects (e.g. car models currently on sale)
- $X$  set of features (e.g. electric, 3 doors, red...)
- $\perp$  incidence relation:  $a \perp x$  iff object  $a$  has feature  $x$
- $R \subseteq A \times X$  knowledge/perception/beliefs of a given agent:  
 $aRx$  iff object  $a$  has feature  $x$  according to the agent
- $a^\uparrow$  set of features of object  $a$
- $x^\downarrow$  set of objects having feature  $x$
- $B^\uparrow$  set of features shared by all objects in  $B$
- $Y^\downarrow$  set of objects satisfying all features in  $Y$
- $\mathbb{P}^+$  concept lattice arising from database  $\mathbb{P}$

# Categories as social constructs

**Social interaction** is key to categorization theory:

- categories arise from factual information about the world.
- **However**, what they mean critically depends on how people **perceive** them and **agree** about them

**Three aspects of categorization theory:**

- factual truth
- subjective perception / knowledge / beliefs
- social interaction

# Epistemic interpretation of $\Box$

In an RS-frame  $\mathbb{F} = (\mathbb{P}, R)$ :

- $R \subseteq A \times X$  encodes perception of a given agent about objects and their features
- $aRx$  reads 'object  $a$  has feature  $x$  according to the agent'
- $\Box\phi$  reads 'category which the agent understands as  $\phi$ '

**Factivity of knowledge.**  $\Box\phi \leq \phi$

$$\forall p(\Box p \leq p)$$

$$\text{iff } \forall \mathbf{m}(\Box \mathbf{m} \leq \mathbf{m})$$

$$\text{iff } \forall a \forall \mathbf{m}[\text{ST}_a(\Box \mathbf{m}) \rightarrow \text{ST}_a(\mathbf{m})]$$

$$\text{iff } \forall a \forall m(aRm \rightarrow a \perp m),$$

if  $a$  has  $m$  according to the agent, then  $a$  has  $m$  in reality

Epistemic interpretation of  $\Box$ , continued

**Positive introspection.**  $\Box\phi \leq \Box\Box\phi$

$$\begin{aligned} & \forall p(\Box p \leq \Box\Box p) \\ \text{iff} & \quad \forall \mathbf{m}(\Box \mathbf{m} \leq \Box\Box \mathbf{m}) \\ \text{iff} & \quad \forall a \forall \mathbf{m}(ST_a(\Box \mathbf{m}) \rightarrow ST_a(\Box\Box \mathbf{m})) \\ \text{iff} & \quad \forall a \forall m(aRm \rightarrow R^{-1}[m]^\uparrow \subseteq R[a]), \end{aligned}$$

**Remark.** Recall that for every  $(B, Y) \in \mathbb{P}^+$ ,

$$\forall a \forall y[(a \in B \ \& \ y \in Y) \rightarrow a \perp y].$$

**Meaning of the latter statement.**

By definition,  $\Box \mathbf{m} := (R^{-1}[m], R^{-1}[m]^\uparrow) \in \mathbb{P}^+$

$$\forall a \forall y[(a \in R^{-1}[m] \ \& \ y \in R^{-1}[m]^\uparrow) \rightarrow aRy].$$

# Social interaction

**RS-frame.** Two agents  $(\mathbb{P}, R_1, R_2)$ .

**Axioms.**  $\Box_i p \leq p$  and  $\Box_i p \leq \Box_i \Box_i p$  for  $1 \leq i \leq 2$ .

**Common knowledge type operator.** For every  $u \in \mathbb{P}^+$ , let

$$C(u) := \bigwedge_{s \in S} su,$$

where  $s \in S$  is such that  $s = (\Box_i \Box_j)^n$  or  $s = \Box_i (\Box_j \Box_i)^n$  with  $i, j \in \{1, 2\}$ ,  $i \neq j$  and  $n \in \mathbb{N}$ .

**Lemma.**  $C$  is a fixed point operator.

**Intuition.** Fixed points of  $C$  are understood as socially constructed categories, the membership and description of which are socially agreed upon.

# Conclusions & further research

## Conclusions

- Categorization theory as a conceptual framework of reference to make sense of RS-semantics for lattice-based modal logic
- Epistemic interpretation capturing knowledge **directly**, rather than by way of uncertainty
- Factivity, positive introspection **retain** their familiar meaning
- Mathematical framework to study categorization phenomena in social science

## Further research

- lattice-based fixed point logics as an epistemic framework for categories socially agreed upon (default categories)
- Formalizing dynamics of categories
- Logical environment for analyzing decision-making procedures